

Mind Reading and Universal Cycles

Does the presenter really possess the power of extrasensory perception?

1. Card Trick

The performer has a deck of 31 cards. Each of 5 audience members removes cards from the case, drops the case on the floor and then gives the deck a straight cut at a random position. The last spectator takes the top card. The deck is then passed back, so each of the 5 spectators take the current top card. The performer now asks, "This may sound strange, but would each of you please look at your card, make a mental picture, and try to send it to me telepathically?" At this is done the performer concentrates and appears confused: "You're doing a great job, but there is too much information coming in for me to make sense of. Would all of you who have a red card please stand up and concentrate?" Now the performer knows all cards.

Possible order of the cards: (31 cards, A, ..., 8 of each suit, except of 8C)

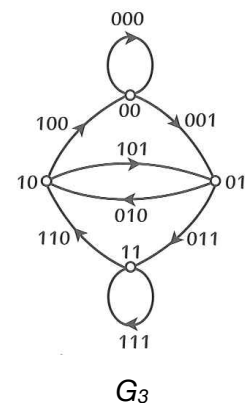
- Write each card in binary number with 5 digits
- First two digits represent the suit (C = 00, S = 01, D = 10, H = 11)
- last three digits are the value in binary numbers mod 8 (i.e. A = 001, 2 = 010, ... , 7 = 111, 8 = 000)
- e.g. 4H = 11100, 8S = 01000, 00010 = 2C
- the received color pattern represent the first card
- the next card can be computed by taking the last 4 digits of the card before as the first 4 and adding the first and the third digit of the card before mod 2 to receive the 5th digit (e.g. 6S = 01110 -> next card: 11101 = 5H)

2. Universal Cycles

2.1. De Bruijn cycles

A *De Bruijn cycle* with window length k is a cyclic zero/one sequence of length 2^k such that every k consecutive digits appear just once, e.g. 00111010 is a De Bruijn sequence of length 3

De Bruijn Graph G_k : Graph with the vertices being the strings of zero/ one symbols of length $k-1$. An edge is going from one vertex x to another vertex y if there is a zero/one length of length k that has x at its left and y as its right. A Eulerian Cycle in G_k gives a De Bruijn Cycle of window length k .



Number of De Bruijn Cycles of window length k : $2^{2^{k-1}-k}$

Construction (Greedy-algorithm)

Given m symbols e_1, \dots, e_m .

- I. Each of the first $k-1$ symbols is chosen equal to e_1 .
- II. The symbol a_m to be added to the sequence (2)

$$a_1 a_2 \dots a_{k-1} a_k \dots a_{m-k+1} \dots a_{m-1}$$

$$a_1 = a_2 = \dots = a_{k-1} = e_1, m \geq k$$

where the a 's stand for the e 's in a certain order, is the e_i with the greatest subscript consistent with the requirement that the section $a_{m-k+1} \dots a_{m-1} a_m$ duplicate no previously occurring section of n symbols in (2).

- III. Rule II is first applied for $m = k$ (in which case $a_m = a_k = e_m$) and is then applied repeatedly until a further application is impossible.

Example:

$k=3, m = 2$ with $e_1 = 0, e_2 = 1$

001110100 => De Bruijn Cycle: 00111010

Applications of de Bruijn Cycles

- Robotic vision

- Cryptography
- DNA sequencing

2.2.U-Cycles in general

Let \mathcal{F}_n be a Family of combinatorial objects of 'rank n ' with their number $l := |\mathcal{F}_n|$. Each $F \in \mathcal{F}_n$ is specified by a sequence $\langle x_1, \dots, x_n \rangle$, where $x_i \in A$, for some fixed alphabet A .

$U = (a_0, a_1, \dots, a_{m-1})$ is a *universal cycle* for \mathcal{F}_n (or *U-cycle* for short) if

$\langle a_{i+1}, \dots, a_{i+n} \rangle$, $0 \leq i < m$, runs through each element of \mathcal{F}_n exactly once, where index addition is performed modulo n .

In the case of De Bruijn cycles the Family is defined as the following:

$$\mathcal{F}_n = B_n = \{0, 1\}^n = \{(x_1, \dots, x_n) \mid x_i \in \{0, 1\}, 1 \leq i \leq n\}, m=2^n$$

2.3.Other Universal Cycles

2.3.1.Permutations

Let S_n be the set of all $n!$ permutations of $\{1, 2, \dots, n\}$. If $\bar{a} = (a_1, a_2, \dots, a_n)$ and $\bar{b} = (b_1, b_2, \dots, b_n)$ are n -tuples of distinct integers we will say \bar{a} and \bar{b} are *order-isomorphic* if

$$a_i < a_j \Leftrightarrow b_i < b_j$$

e.g. for $n=5$: (1, 3, 2) and (2, 5, 3) are order-isomorphic.

A U-Cycle $U_n = (a_0, a_1, \dots, a_{n!-1})$, $a_i \in \{1, 2, \dots, N\}$ for S_n will be a $n!$ -tuple such that each $\sigma \in S_n$ is order-isomorphic to exactly one block $(a_{i+1}, a_{i+2}, \dots, a_{i+n})$. (index addition is performed modulo $n!$)

e.g. the elements of S_3 are 123, 132, 213, 231, 312, 321, so a U-Cycle for S_3 is 145243

2.3.2.Partitions

Let P_n be the set of partitions of the n -element set $\{1, 2, \dots, n\}$.

A U-Cycle will be a sequence composed of symbols from the set $A = \{a, b, c, \dots\}$ containing every partition exactly once. A block represents a partition, in which i and j are in the same group if the i th and the j th symbol are the same (e.g. for $n=5$ abacc represents the partitions 13|2|45).

e.g. there are 52 partitions of $\{1, 2, 3, 4, 5\}$, a U-Cycle on the alphabet $\{C, S, D, H, J\}$ is DDDDDCHHHCCDDCCCHCHCSHHSDDSSSHSDDCHSSCHSHDHSHSJSJDCD

The existence of such U-cycles depends on n . For example, there is no U-cycle for $n=3$.

2.3.3.k-sets of a n-set

$\left[\begin{smallmatrix} n \\ k \end{smallmatrix} \right] :=$ family of all k -element subsets of an n -element set $\{0, 1, \dots, n-1\}$.

U-cycle: ($k=3, n=8$):

02456145712361246703671345034601250135672560234723570147.

Each subsequence abc represents $\{a, b, c\}$

Existence: If k does not divide $\binom{n-1}{k-1}$, $\left[\begin{smallmatrix} n \\ k \end{smallmatrix} \right]$ does not divide, has no U-cycle

3. References

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